

Lab 10: Propositions as types

Functional Programming (ITI0212)

This week we are learning about the *propositions as types* paradigm. It tells us that propositions (the mathematical statements) are in fact the same things as types (the functional programming objects). This is known as the Curry-Howard correspondence. When we have $t : T$, then it can be read both as *t is a term of type T* and *t is a proof of the proposition T*. We see here that terms becomes proofs, and constructing a term of a type (= proposition) is the exact same thing as constructing a proof of this type (= proposition).

Facts about less or equal

As a general hint for this section, do not forget that you can use previous result to prove new ones. Recall the type \leq' from the lecture:

Task 1

Prove that \leq' is transitive, which means you need to give a function of the following type: $\text{leTrans} : \{m\ n\ p : \text{Nat}\} \rightarrow (m \leq' n) \rightarrow (n \leq' p) \rightarrow (m \leq' p)$

Task 2

State and prove that any integer is less or equal than its successor. This will be the function of the following type: $\text{succLarger} : \{n : \text{Nat}\} \rightarrow n \leq' (n + 1)$

Task 3

Prove the following: $\text{leWeakRight} : \{m\ n : \text{Nat}\} \rightarrow (m \leq' n) \rightarrow m \leq' (n + 1)$ and $\text{leWeakLeft} : \{m\ n : \text{Nat}\} \rightarrow ((m + 1) \leq' n) \rightarrow (m \leq' n)$

Task 4

Prove the following: $\text{zeroPlusLeft} : \{m\ n : \text{Nat}\} \rightarrow (0 + n) \leq' m + n$

On the length of lists

Task 5

State and prove that if two lists xs , ys are such that $\text{length } xs \leq' \text{length } ys$, then $\text{length } (x :: xs) \leq' \text{length } (y :: ys)$.